

Compositional Performance Analysis

CPEN 432 Real-Time System Design

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So far ...

- Arrival model

- Aperiodic
- Periodic (with and without offsets)
- Sporadic

- Response time analyses

- Uniprocessor: $R_k = C_k + \sum_{\tau_i \in hp(k)} \left(\left\lceil \frac{R_k}{T_i} \right\rceil \cdot C_i \right)$

- Symmetric multiprocessors: $R_k = C_k + \frac{1}{m} \sum_{\tau_i \in hp(k)} \left(\left\lceil \frac{R_k}{T_i} \right\rceil \cdot C_i + C_i \right)$

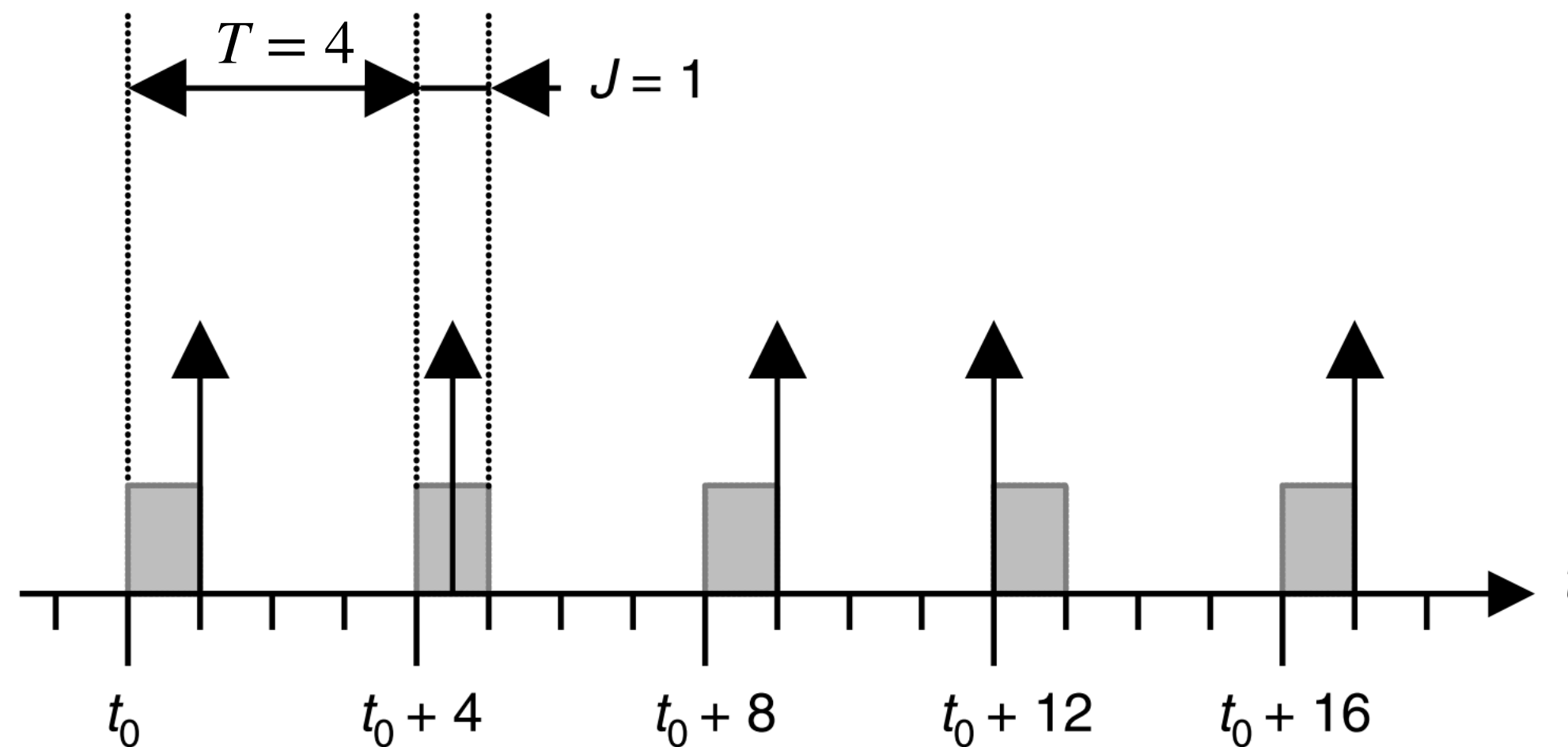
- ...

- In a distributed systems, complex embedded applications, MPSoCs, etc.

- Periodic task model too rigid, as it does not account for variations

“Periodic with Jitter” Event Model

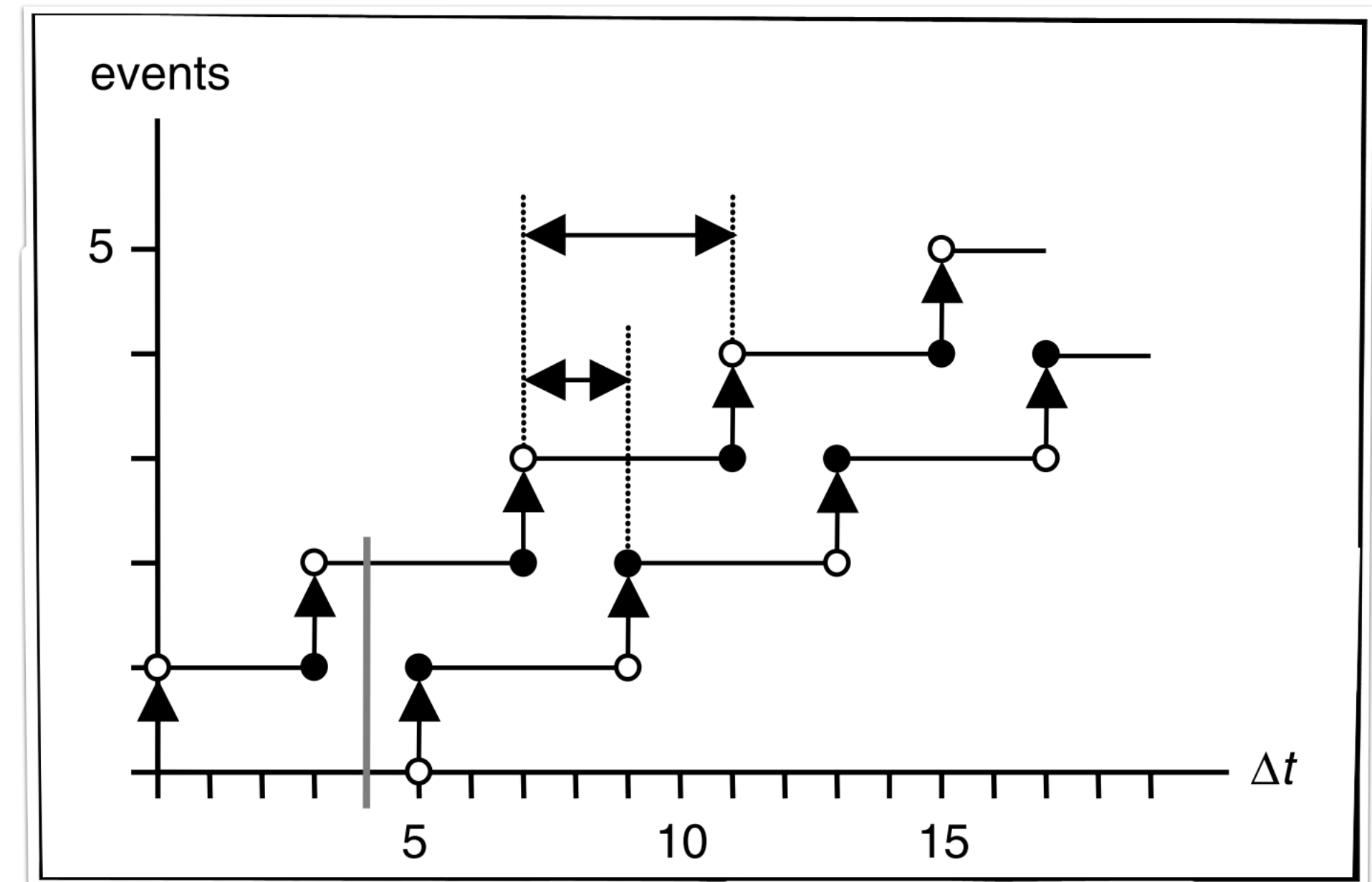
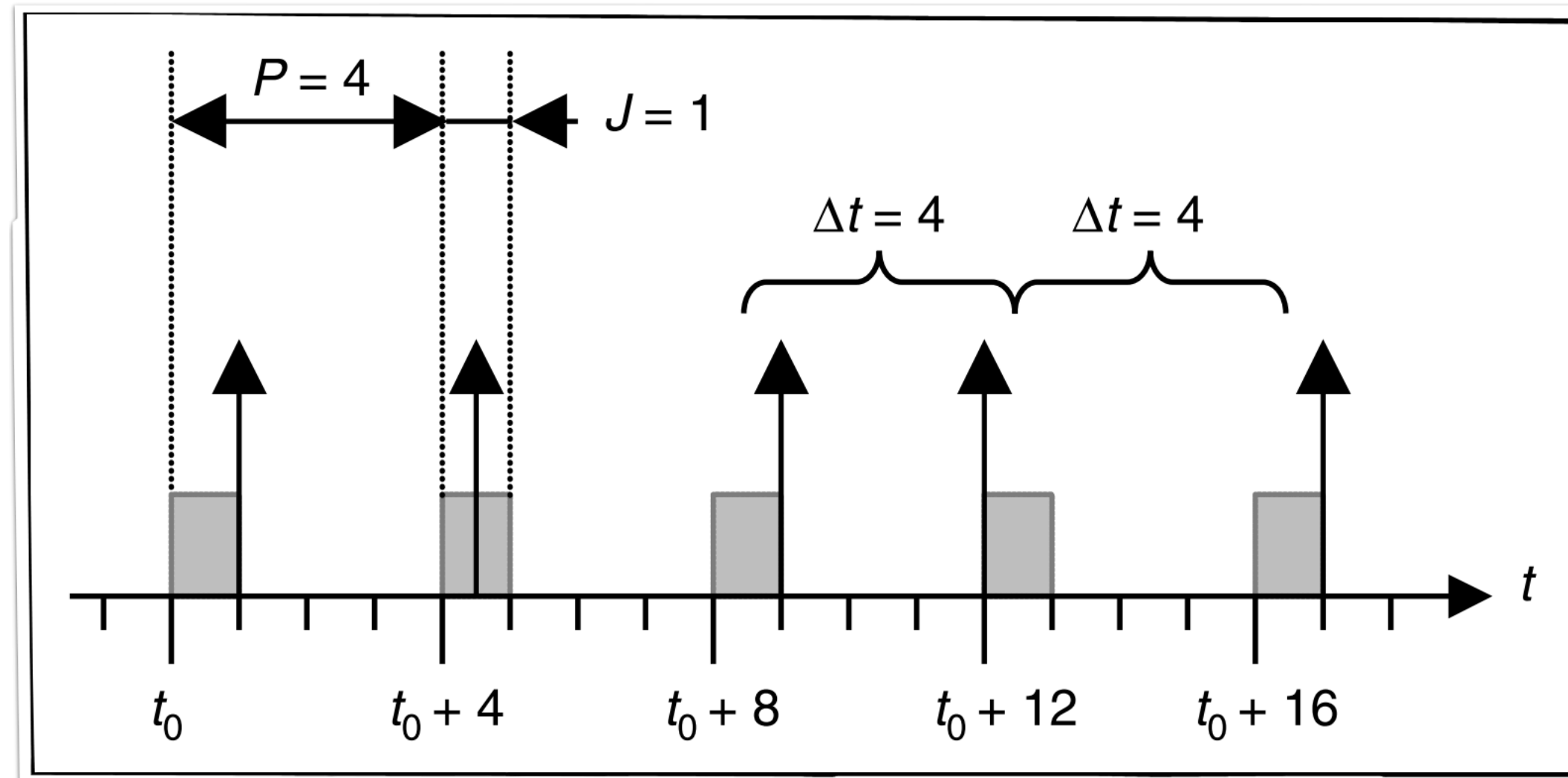
- Two parameters: period T and jitter J
 - Each event **generally occurs periodically** with period T
 - But it **can vary a bit** around its exact position within a jitter interval J
- Example event stream satisfying $(T, J) = (4, 1)$



Event Functions

Definition 1 (Upper event function): The upper event function $\eta^u(\Delta t)$ specifies the maximum number of events that can occur during any time interval of length Δt .

Definition 2 (Lower event function): The lower event function $\eta^l(\Delta t)$ specifies the minimum number of events that have to occur during any time interval of length Δt .



$$\eta_{T+J}^u(\Delta t) = \left\lceil \frac{\Delta t + J}{T} \right\rceil$$

$$\eta_{T+J}^l(\Delta t) = \max \left(0, \left\lfloor \frac{\Delta t - J}{T} \right\rfloor \right)$$

Distance Functions

Definition 3 (Minimum distance function): The minimum distance function $\delta^{\min}(N \geq 2)$ specifies the minimum distance between $N \geq 2$ consecutive events in an event stream.

Definition 4 (Maximum distance function): The maximum distance function $\delta^{\max}(N \geq 2)$ specifies the maximum distance between $N \geq 2$ consecutive events in an event stream.

$$\delta^{\min}(N \geq 2) = \max\{0, (N - 1) * T - J\}$$

$$\delta^{\max}(N \geq 2) = (N - 1) * T + J$$

- Questions

- How can you define $\eta^u(\Delta t)$ and $\eta^l(\Delta t)$ in terms of $\delta^{\min}(N)$ and $\delta^{\max}(N)$?

- $\eta^u(\Delta t) = \max_{n \geq 1, n \in \mathbb{N}} \{n \mid \delta^{\min}(n) < \Delta t\}$

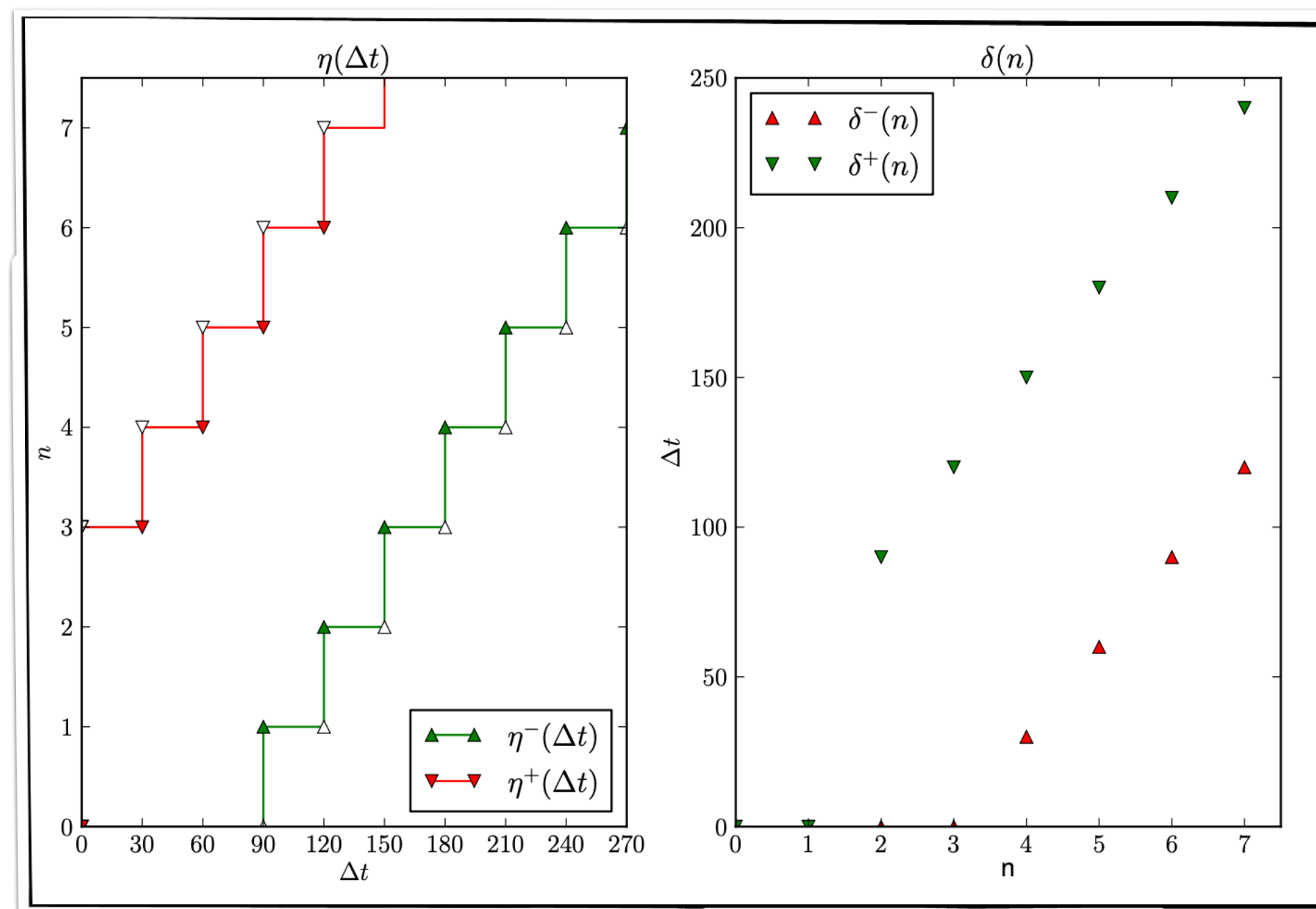
- $\eta^l(\Delta t) = \min_{n \geq 1, n \in \mathbb{N}} \{n \mid \delta^{\max}(n + 2) > \Delta t\}$

Questions

- For sporadic events with a minimum inter-arrival time T
 - Is the jitter parameter J meaningful?
 - How is the lower event function $n^l(\Delta t)$ defined?
 - How is the maximum distance function $\delta^{max}(N \geq 2)$ defined?

Questions

- What happens if $J > T$ in a “periodic with jitter” event model?
 - Two or more events can occur at the same time, leading to **bursts**
 - New parameter d_{min} that captures the minimum distance between events in a burst
 - Can you draw event and distance functions for $(T, P) = (30, 60)$?



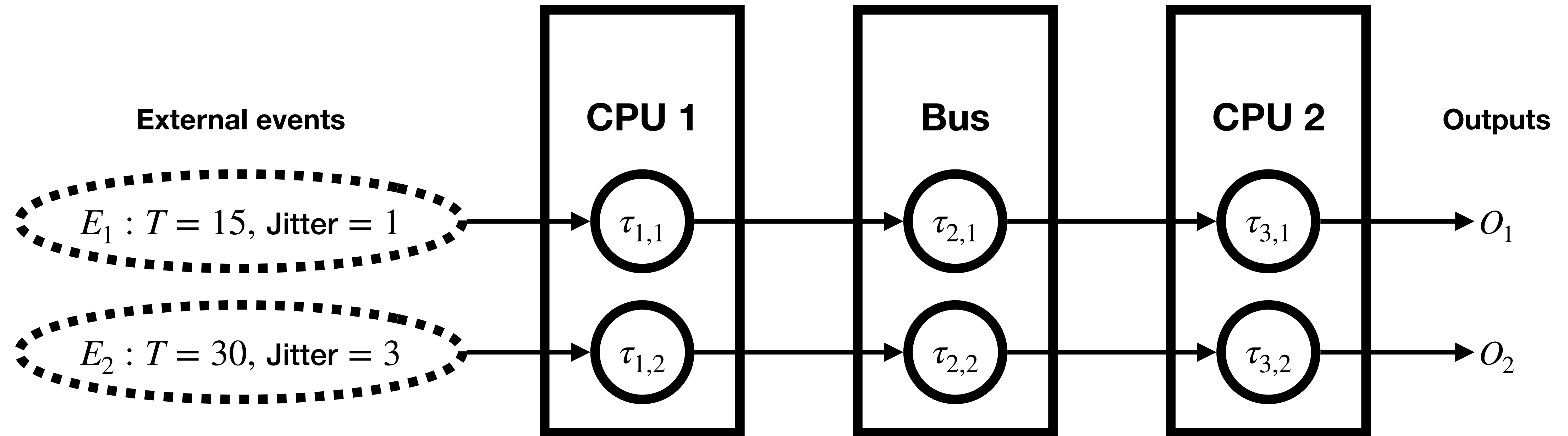
Response-Time Analysis (Local)

- Compute the maximum q -event busy window $B_i^+(q)$
 - An upper bound on the time a resource requires to service q activations of task τ_i
 - Assumption: all q activations arrive “sufficiently early”
 - i.e., q^{th} event arrives prior to the completion of its preceding event (the $(q - 1)$ -event busy window)
- For fixed-priority preemptive scheduling,
 - Starting with $B_k^u(q) = q \cdot C_k$, solve for $B_k^u(q) = q \cdot C_k + \sum_{i \in hp(k)} \eta_i^u(B_k^u(q)) \cdot C_i$
 - Stopping condition
 - Consider only the first q_k^u activations, where $q_k^u = \min\{q \in \mathbb{N}^+ \mid \delta_k^{min}(q + 1) \leq B_k^u(q)\}$
 - $R_k^u = \max_{q \in \mathbb{N}^+ \mid q \leq q_k^u} (B_k^u(q) - \delta_i^{min}(q))$

Output Event Function of a Task

- R_k^u from response-time analysis, $R_k^l = C_k$
 - Thus, the scheduling policy adds an additional jitter of $R_k^u - R_k^l$
 - That is, the output jitter is $J_{k,out} = J_k + (R_k^u - R_k^l)$
 - Often, J_k is denoted as $J_{k,in}$ or $J_{k,act}$
- The output event model period obviously equals the activation period
 - That is, $P_{k,out} = P_{k,in}$

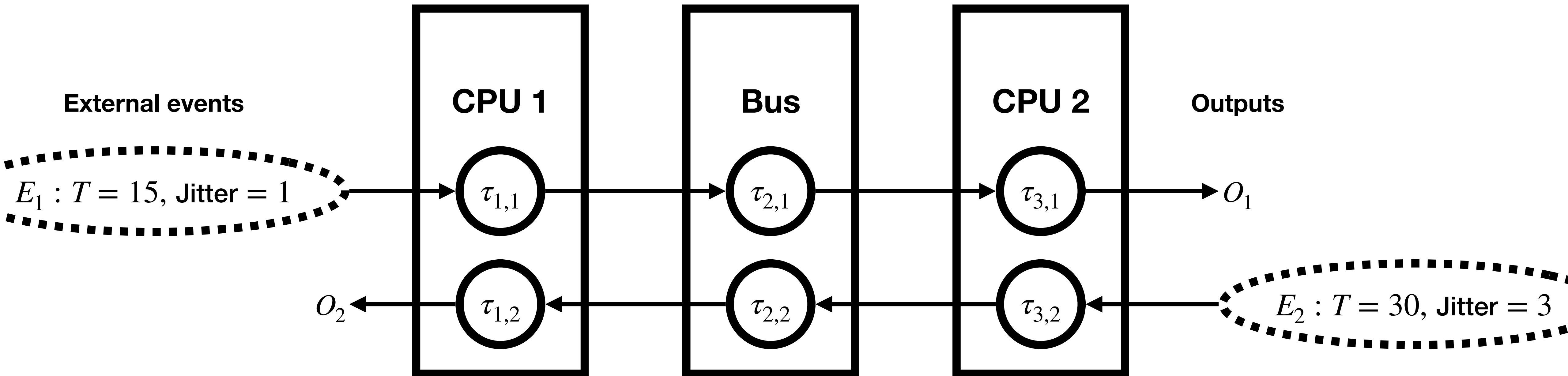
Example 1



What is the end-to-end path latency of $E_1 \rightarrow O_1$?

What is the end-to-end path latency of $E_2 \rightarrow O_2$?

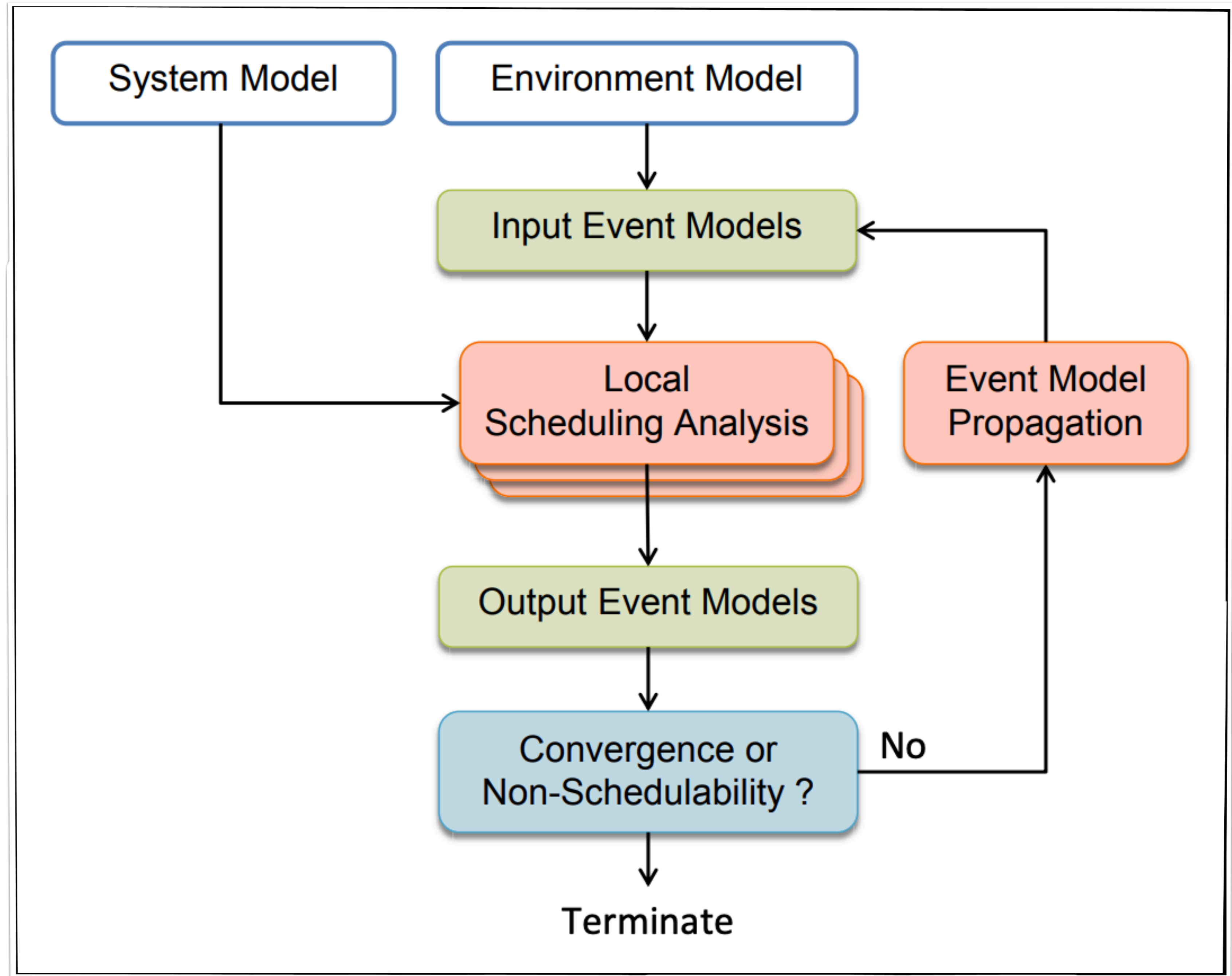
Example 2



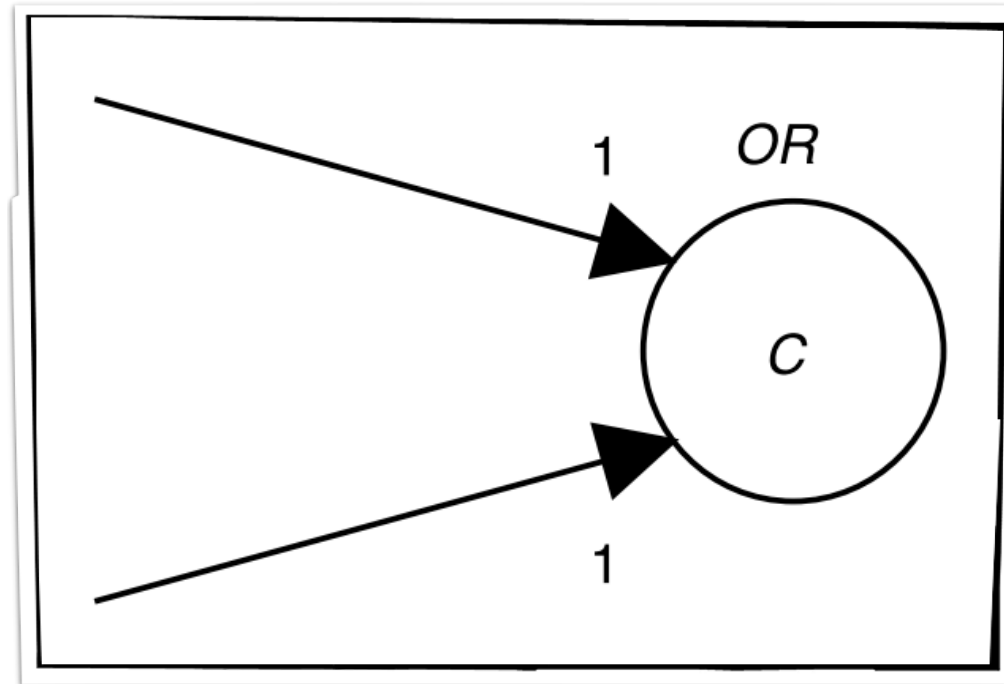
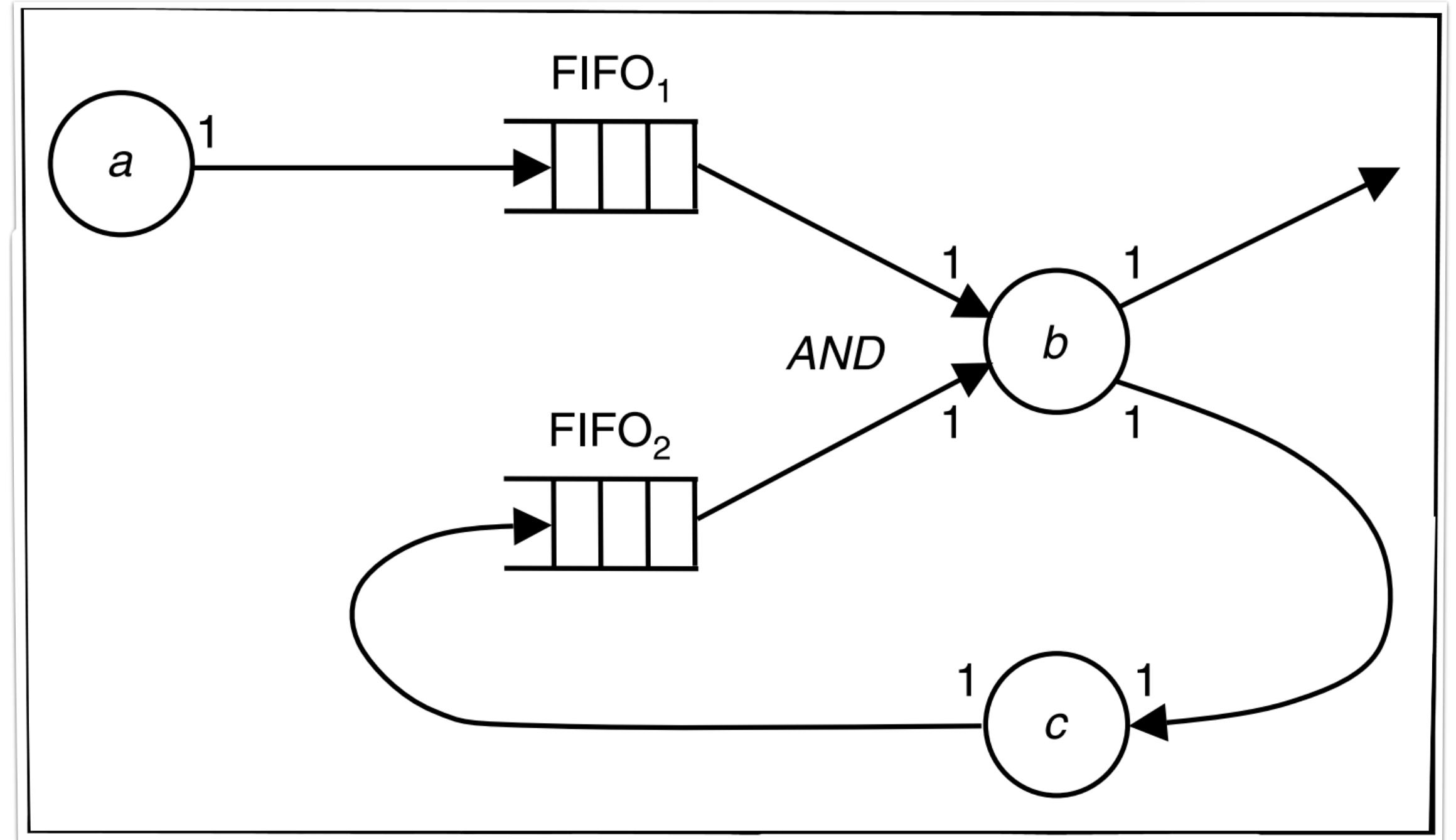
What is the end-to-end path latency of $E_1 \rightarrow O_1$?

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Compositional Performance Analysis

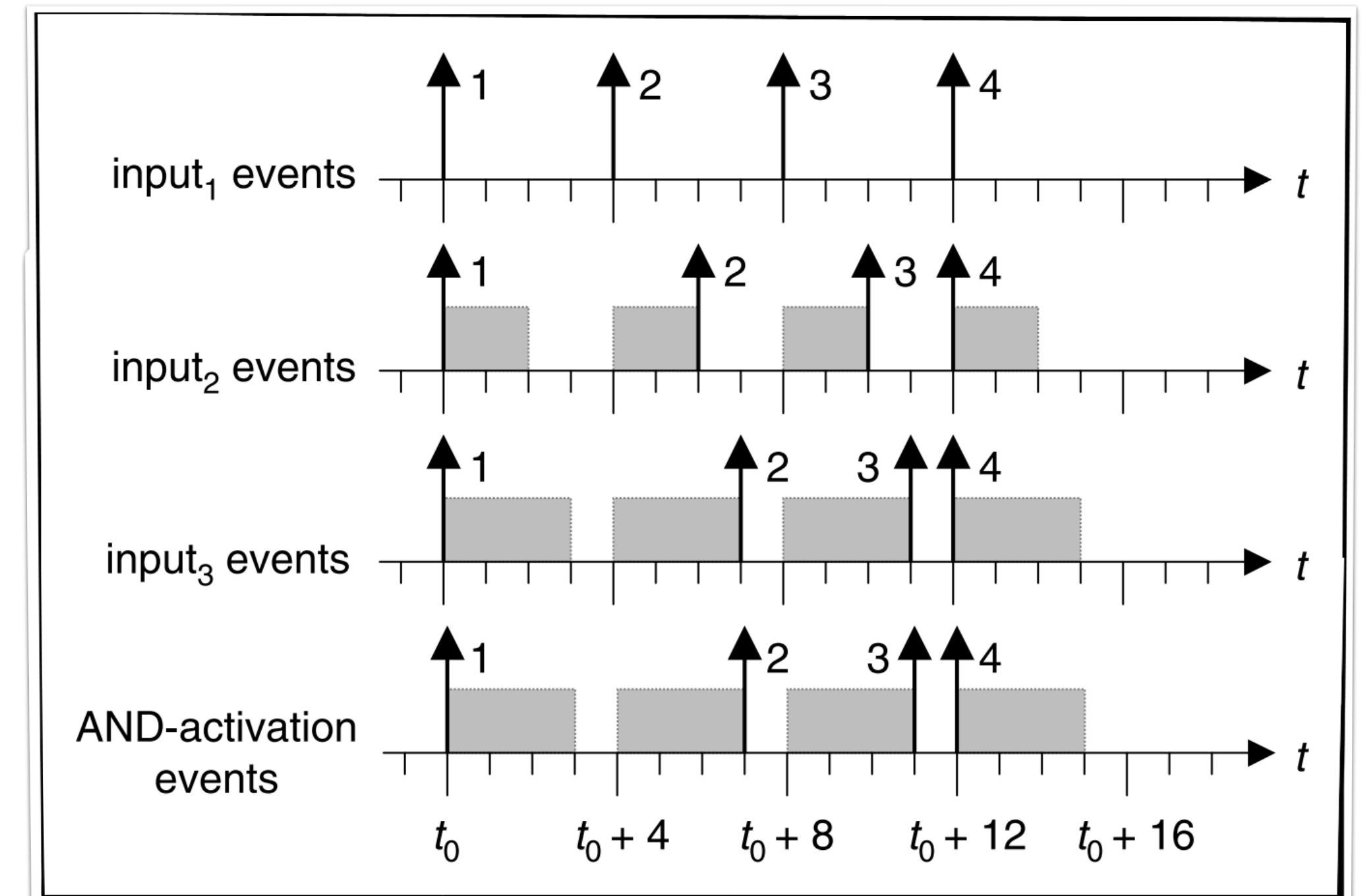
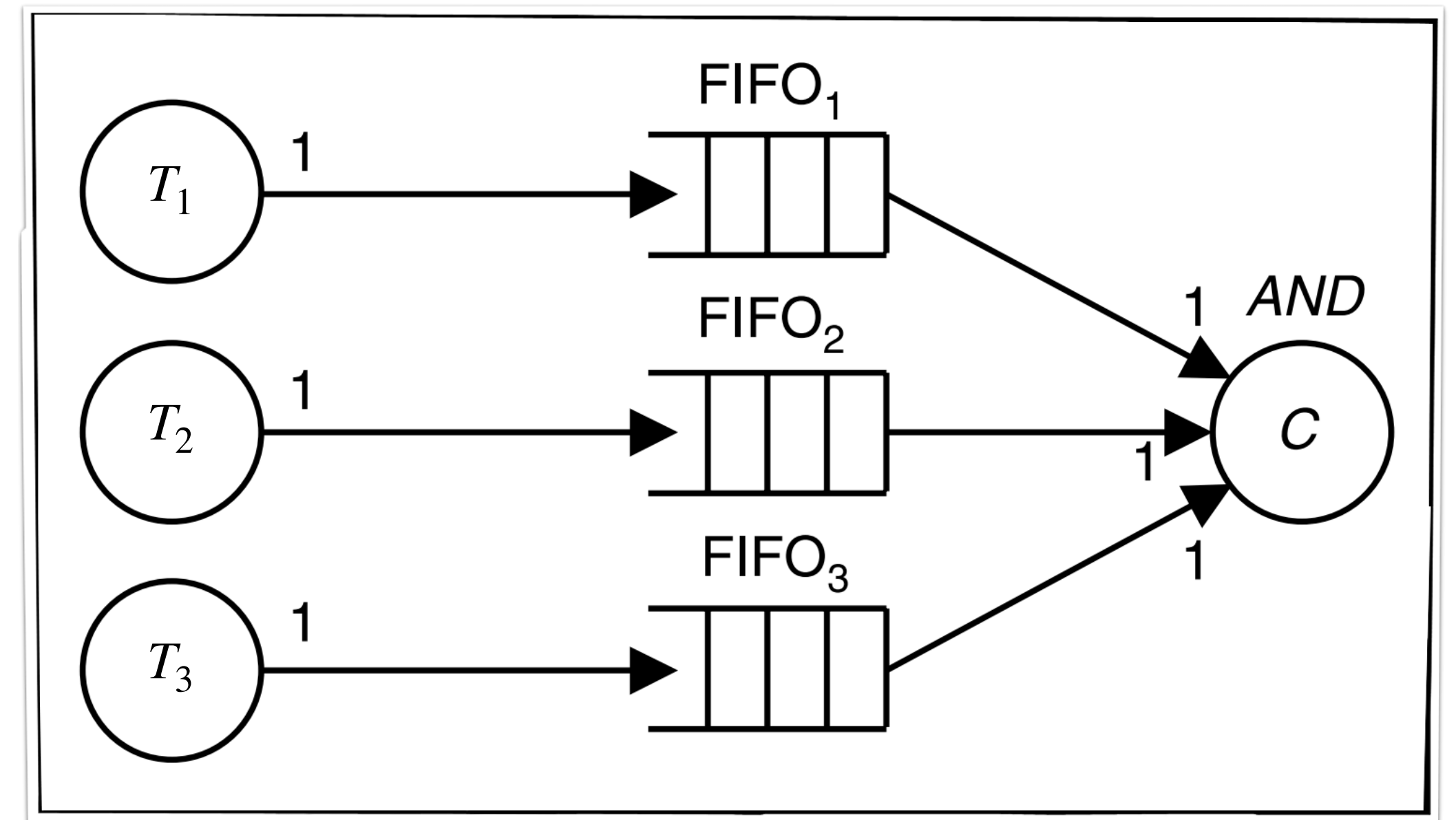


The diagram illustrates a parallel processing system. On the left, three circular nodes labeled T_1 , T_2 , and T_3 are arranged vertically. Each node has an outgoing arrow pointing to a corresponding FIFO buffer. The arrows are labeled with the value '1'. The FIFO buffers are labeled $FIFO_1$, $FIFO_2$, and $FIFO_3$ from top to bottom. Each FIFO buffer is represented as a horizontal rectangle divided into four vertical slots. From the right side of each FIFO buffer, an arrow points to a circular node labeled C . These arrows are also labeled with the value '1'. Above the node C , the word 'AND' is written, indicating that node C is an AND gate that receives three inputs, each with a value of 1.

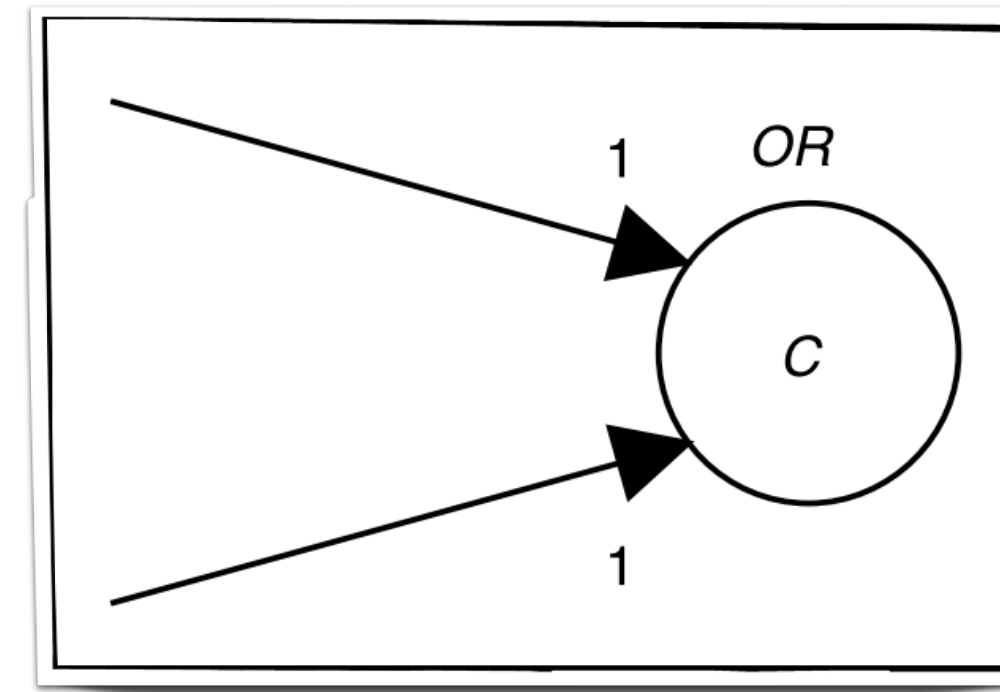


AND-activation

- FIFO channels?
 - Input data buffering: Data may have to wait at some inputs until all other inputs have the necessary data
- Tokens?
 - Amount of data required per input event
- AND-activation period?
 - “To ensure bounded AND-buffer sizes the period of all input event models must be the same.”
 - Example: $T_{AND} = T_1 = T_2 = T_3$
- AND-activation jitter?
 - Example: $J_{AND} = \max\{J_1, J_2, J_3\}$

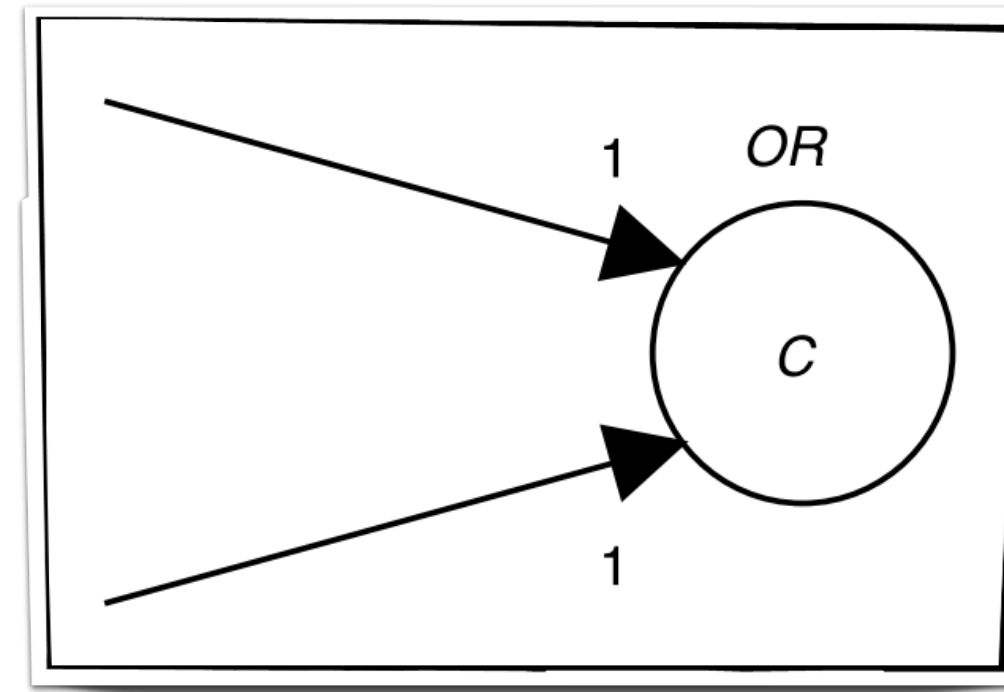


OR-activation



- No FIFO channels. Why?
 - Data at one input never has to wait for data to arrive at a different input for activation
- OR-activation period?
 - Example input event models
 - Event stream 1 : $T_1 = 4, J_1 = 2$
 - Event stream 2 : $T_2 = 3, J_2 = 2$

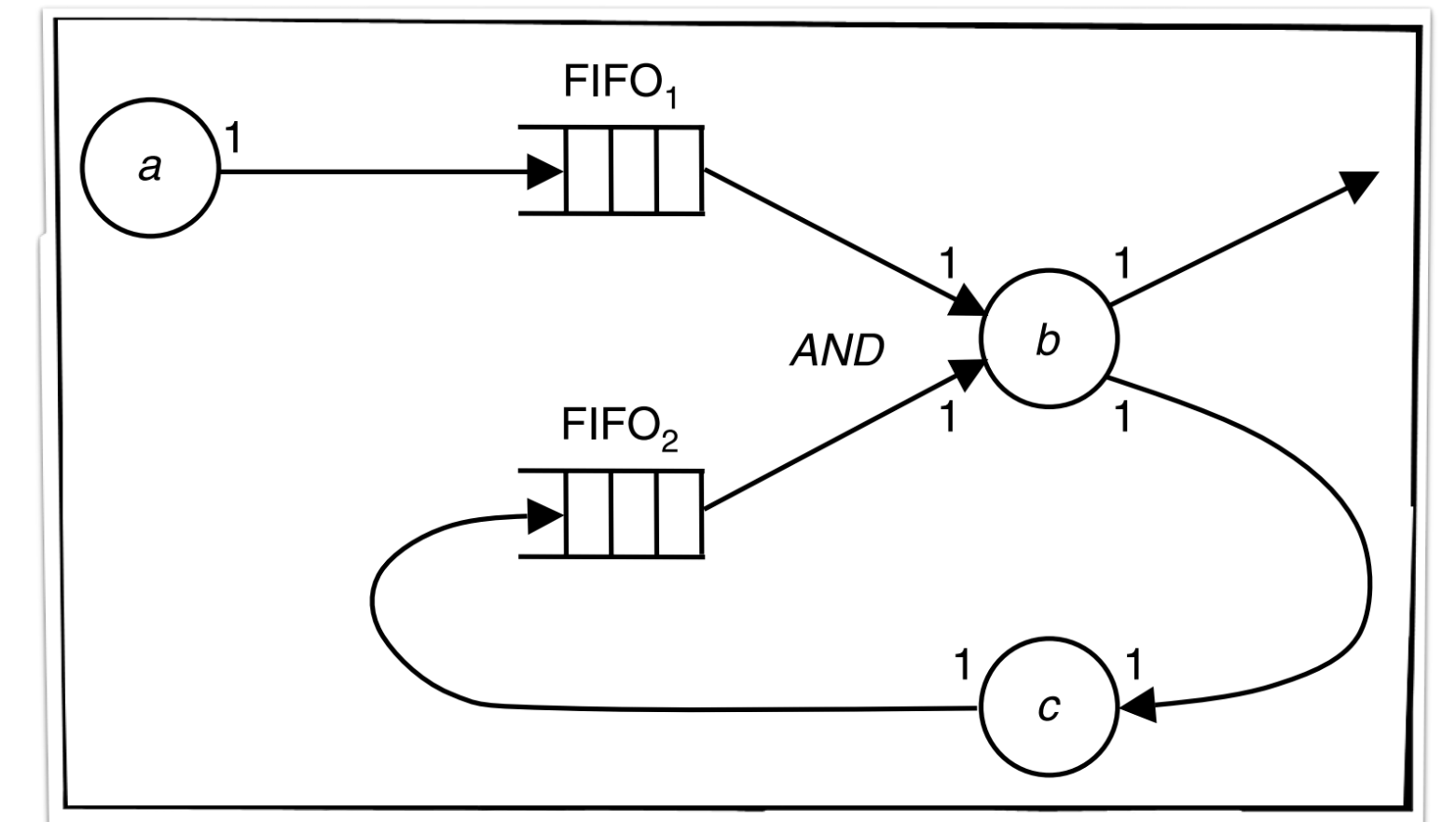
OR-activation



- No FIFO channels. Why?
 - Data at one input never has to wait for data to arrive at a different input for activation
- OR-activation period?
 - Example input event models
 - Event stream 1 : $T_1 = 4, J_1 = 2$
 - Event stream 2 : $T_2 = 3, J_2 = 2$
 - Consider the *macro* period — the least common multiple of all input event model periods
 - $T_{macro} = LCM(T_i)$
 - Consider the total events across all input streams arriving in the macro period (assuming zero jitter)
 - $$N_{total} = \sum_{i=1}^n \frac{LCM(T_i)}{T_i}$$
 - The OR-activation period is thus the average period, given by
 - $$T_{OR} = \frac{T_{macro}}{N_{total}} = \frac{1}{\sum_{i=1}^n \frac{1}{T_i}}$$

Cyclic Task Dependencies

- Idea
 - Recall that AND-activation jitter is the maximum of input jitters
 - Start with zero jitter for the cyclic input
 - Update, if there are changes
 - This approach may not work. Why?
- Problem
 - Timing of cycle-external and cycle-internal inputs is correlated
 - But the AND-activation jitter (maximum of input jitter) ignores this
- Solution
 - Consider possible phases between inputs arriving at cycle-external and cycle-internal inputs

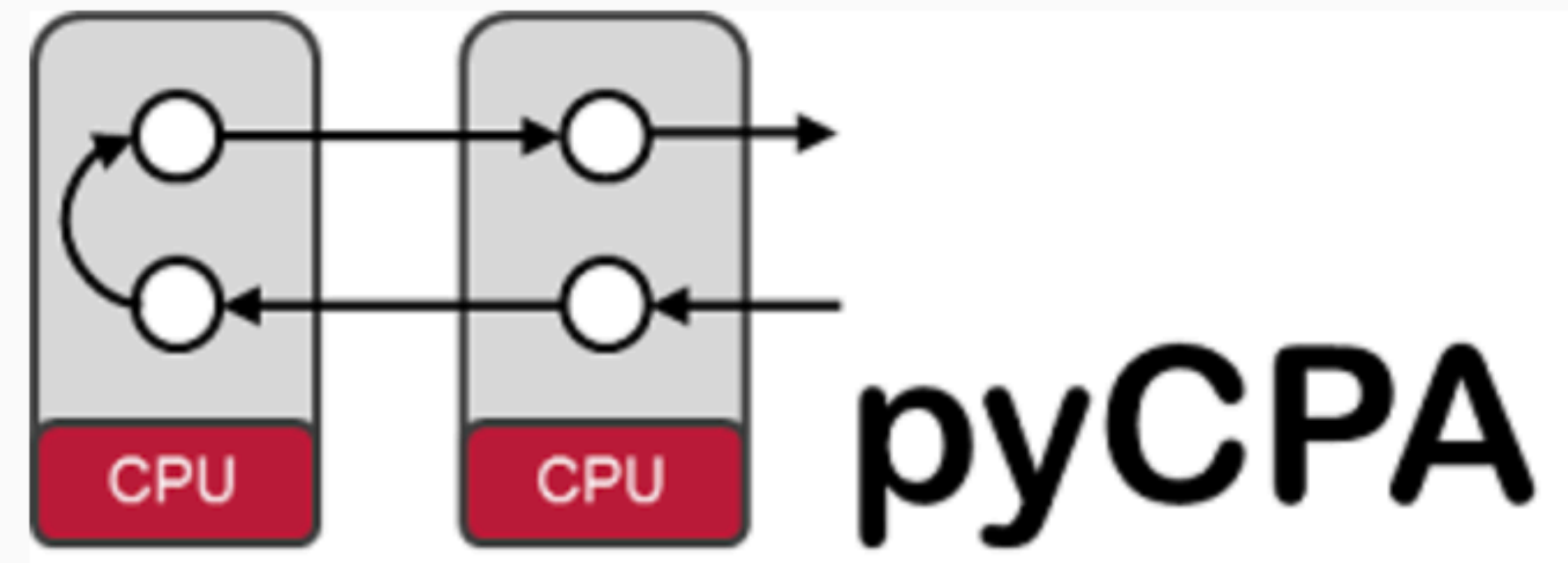


Open-Source Tool

<https://github.com/IDA-TUBS/pycpa> | <https://pycpa.readthedocs.io/>

Welcome

pyCPA is a pragmatic Python implementation of Compositional Performance Analysis (aka the SymTA/S approach provided by [Symtavisision](#) (now: [Luxoft](#))) used for research in worst-case timing analysis. Unlike the commercial SymTA/S tool, pyCPA is not intended for commercial-grade use and does not guarantee correctness of the implementation.



Thank You!