

Variable Unification

Unification is Prolog's matching technique.

1 Unbound variables unify with anything.

2 Constants or integers only unify with itself.

3 Structures only unify with other structures if: (a) same name and number of arguments, (b) all the corresponding arguments unify.

1) $100=10*10 \rightarrow$ No, case 2 ("is" would work)

2) $\text{struct}(A, b(C, d), e) = \text{struct}(X, X, Y) \rightarrow$ Yes, case 3

3) $[a,b,c] = [a|[b,c]] = [a,b|[c]] = [a,b,c|[]] \rightarrow$ Yes, case 3

4) $[1|[2, 3]] = [\text{Last}|\text{First}] \rightarrow$ Yes, Last = 1, First = [2, 3]

5) $[1, Y, 3] = [A|B], 25 = Y. \rightarrow$ Yes

6) $[[[a, b]] | c] = [H|T] \rightarrow$ Yes, $H = [[a, b]]$, $T = c$

7) $[a(25, b), c(B), B|T] = [X, c(400), D, D] \rightarrow$ Yes, $B = 400$, $T = [400]$, $X = a(25, b)$, $D = 400$

8) $\text{oh}(MY) = MY \rightarrow$ infinite recursion, unifies as $MY = \text{oh}(**)$

9) $\text{Ans} = 25*4 \rightarrow$ Yes, $\text{Ans} = 25*4$

10) $[.] = [X|Y] \rightarrow$ Yes, $X = .$, $Y = []$

11) $Y \text{ is } (6+6)/3 \rightarrow$ Yes, $Y = 4.0$

12) $6+6 \text{ is } 3*4 \rightarrow$ No, $6+6$ is a structure

13) $\text{Tallgeese is mobilesuit}(0) \rightarrow$ Error, "is" means it tries to evaluate mobilesuit as a mathematical operation

14) $\text{assert}(\text{match}(\text{stick})), \text{retract}(\text{match}(X)) \rightarrow X=\text{stick}$

15) $\text{assert}(\text{match}(\text{stick})), \text{retract}(\text{match}(X)), \text{retract}(\text{match}(X)) \rightarrow$ False, last retract fails to find a match, whole query fails

List Manipulation

?- $\text{count}([1,2,3],N)$. $N = 3$

$\text{count}([],0)$. % initializes Sum to 0 (when list has been traversed)

$\text{count}([_|T],N) :-$

$\text{count}(T,\text{Sum})$, % recursively traverse the list

$N \text{ is } \text{Sum}+1$. % increase counter during backtracking

?- $\text{tally}([a,b,b,c],b,N)$. $N = 2$

$\text{tally}(L,Q,_)$:-

$\text{count}(L,Q,0)$. % call count with N initialized to 0

$\text{count}([Q|T],Q,N) :-$ % if Q matches current element (i.e. head)

!, % cut to not search exhaustively

$N1 \text{ is } N+1$, % increase counter

$\text{count}(T,Q,N1)$. % recursively traverse the list with the new count

$\text{count}([_|T],Q,N) :-$ % else Q does not match current element (because above rule already checked for match)

%not($H = Q$) % can add a line such as not($H = Q$) or $H \neq Q$ to be safe, but not really needed

$\text{count}(T,Q,N)$. % recursively traverse the list with the same count

$\text{count}([],_,N) :-$ % if finished traversing the list,

$\text{write}('N = ')$,

$\text{write}(N)$, % print out value of N. if this isn't done, then backtracking causes N to go back to "something"

nl .

?- $\text{removeAll}(b,[a,b,b,c],M)$. $M = [a,c]$

$\text{removeAll}(A,[A|Tail],\text{NewList}) :-$

$\text{removeAll}(A,\text{Tail},\text{NewList})$.

$\text{removeAll}(A,[\text{Head}|\text{Tail}],[\text{Head}|\text{NewTail}]) :-$

%not($\text{Head} = A$),

$\text{removeAll}(A,\text{Tail},\text{NewTail})$.

$\text{removeAll}(_,[],[])$.

?- $\text{removeItem}(b,[b,a,b],M)$. $M = [a,b]$; $M = [b,a]$

$\text{removeItem}(A,[A|L],L)$. % base case -- when match found, co

$\text{removeItem}(A,[N|M],[N|L]) :-$ % adds non matching elements

$\text{removeItem}(A,M,L)$.

?- $\text{removeGen}(\text{all},b,[a,b,b],M)$. $M = [a]$

$\text{removeGen}(\text{one},A,L,M) :-$ % remove one item

$\text{removeItem}(A,L,M)$.

$\text{removeGen}(\text{all},A,L,M) :-$ % remove all items

$\text{removeAll}(A,L,M)$.

?- $\text{writeList}([\text{hello},\text{world}])$. 1 hello (nl) 2 world

$\text{writeList}(L) :-$

$\text{writeList}(L,0)$. % need to initialize counter to 0, only

$\text{writeList}([H|T],N) :-$

$N2 \text{ is } N+1$, % increase counter

$\text{write}(N2)$, % write counter

$\text{write}(' ')$, % write white space

$\text{write}(H)$, % write current element

$\text{writeList}(T,N2)$. % recursively traverse the list with the r

?- $\text{reverseItA}([1,2],\text{New})$. $\text{New} = [2,1]$

$\text{reverseItA}(\text{Old},\text{New}) :- \text{flip}(\text{Old},[],\text{New})$.

$\text{flip}([\text{OldHead}|\text{OldTail}],\text{WorkingList},\text{FinalList}) :-$

$\text{append}([\text{OldHead}],\text{WorkingList},\text{NewWorkingList})$,

$\text{flip}(\text{OldTail},\text{NewWorkingList},\text{FinalList})$.

$\text{flip}([],\text{CompleteWorkingList},\text{CompleteWorkingList})$.

%#2 final list as *something*, set to [] at end of tree

$\text{reverseItB}([\text{OldHead}|\text{OldTail}],\text{New}) :-$

$\text{reverseItB}(\text{OldTail},\text{ReturnedNew})$,

$\text{append}(\text{ReturnedNew},[\text{OldHead}],\text{New})$.

$\text{reverseItB}([],[])$.

%#3 identical to #1, but doesn't use append/3 predicate

element (because above rule already checked for match)

$\text{reverseItC}(\text{Old},\text{New}) :- \text{flipC}(\text{Old},[],\text{New})$.

$\text{flipC}([\text{OldHead}|\text{OldTail}],\text{WorkingList},\text{Final}) :-$

$\text{flipC}(\text{OldTail},[\text{OldHead}|\text{WorkingList}],\text{Final})$.

$\text{flipC}([],\text{CompleteWorkingList},\text{CompleteWorkingList})$.

(i.e. list is currently empty)

% #4

$\text{reverse}(\text{List},\text{Reversed}) :- \text{rev}(\text{List},[],\text{Reversed})$.

$\text{rev}([],\text{Rev},\text{Rev})$.

$\text{rev}([A|T],L,\text{Rev}) :- \text{rev}(T,[A|L],\text{Rev})$.

```

% append two lists
join([],L,L).
join([X|R],S,[X|T]):- join(R,S,T).

intersection([],_,[]).
intersection([X|R],Y,[X|Z]):-
    member(X,Y), intersection(R,Y,Z).
intersection([X|R],Y,Z):-
    not(member(X,Y)),
    intersection(R,Y,Z).

?- union([1,2],[2,3],U). U = [1,2,3]
union([], X, X).
union([X|R],Y,Z):- member(X,Y), union(R,Y,Z).
union([X|R],Y,[X|Z]):- not(member(X,Y)), union(R,Y,Z)

member(X,[X|_]).
member(X,[_|T]):- member(X,T).

?- is_palindrome([b,o,b]). true
is_palindrome(L):- reverse(L,L). % reverse is built-in

even([]).
even([_,_|T]):- even(T).

odd([]).
odd([_,_|T]):- odd(T).

Create a list of integers from 0 to N
intlist(N,L):- N > 0, make_intlist(0, N, L).
make_intlist(M,N,[]): M > N. % exit case
make_intlist(M,N,L):- M <= N, M2 is M + 1,
    make_intlist(M2,N,L2), % M+1,...,N
    append([M],L2,L).

item(bag). item(pc). item(pc).
?- setof(Stuff,item(Stuff),Set).
-> Set = [bag, pc].
?- bagof(Stuff,item(Stuff),Bag).
-> Bag = [bag, pc, pc].

delete_all(A,[A|T],T2):- delete_all(A,T,T2).
delete_all(A,[B|T],[B|T2]):- delete_all(A,T,T2).
delete_all(_,L,L).

Swap first and last elements in a list.
?- transfer([a,b,c],X). X = [c,b,a].
transfer([H|T],[H2|T2]):- swap(T,H,H2,T2).
swap([First,Second|Rest],H,H2,[First|T2]):-
    swap([Second|Rest],H,H2,T2).
swap([Last],H,Last,[H]).

Raise each element to a power of itself.
?-raise([1,2,3,4],X). X = [1, 4, 27, 256].
raise([], []).
raise([H|T],[H2|T2]):- raise(T,T2), H2 is H**H.

% Find last element of a list
last(X,[X]). % base case
last(X,[_|L]):- last(X,L).

?- element_at(X,[a,b,c,d,e],3). X = c
element_at(X,[X|_],1). % base case
element_at(X,[_|L],K):- K > 1, K1 is K - 1,
    element_at(X,L,K1).

?- flatten([a,[b],c],X). X = [a, b, c]
flatten(X,[X]) :- \+ is_list(X).
flatten([], []).
flatten([X|Xs],Zs):-
    flatten(X,Y), flatten(Xs,Ys), append(Y,Ys,Zs).

factorial(0,1).
factorial(X,Y) :- X1 is X - 1, factorial(X1,Z),
    Y is Z*X,!.

remove_dups([1,2,2,3],L). L = [1,2,3]
remove_dups([], []).
remove_dups([A|R],S):- % do not add if member
    member(A,R), remove_dups(R,S), !.
remove_dups([A|R],[A|S]):- remove_dups(R,S).
OR
compress([], []).
compress([X],[X]).
compress([X,Xs],Zs):- compress([X|Xs],Zs).
compress([X,Y|Ys],[X|Zs]):- X\=Y, compress([Y|Ys],Zs).

?- is_prime(7). true.
is_prime(2). is_prime(3). % base cases
is_prime(P):- integer(P), P > 3, P mod 2 =\= 0,
    \+has_factor(P,3).
has_factor(N,L):- N mod L =:= 0.
has_factor(N,L):- L*L<N, L2 is L+2, has_factor(N,L2).

?- odd_calc([1,2,3,4,5],Ans). Ans = 3
odd_calc(List,Answer):- odd_add(List,0,Answer).
odd_add([LHead|LTail],Working,Answer):-
    NewWorking is Working+LHead,
    odd_sub(LTail,NewWorking,Answer).
odd_add([],Answer,Answer).
odd_sub([LHead|LTail],Working,Answer):-
    NewWorking is Working-LHead,
    odd_add(LTail,NewWorking,Answer).
odd_sub([],Answer,Answer).

?- duplicate([a,b],X). X = [a,a,b,b]
duplicate([], []).
duplicate([X|Xs],[X,X|Ys]):- duplicate(Xs,Ys).

?- pow(2,3,X). X = 8
pow(_,0,1).
pow(X,Y,Z) :- Y1 is Y-1, pow(X,Y1,Z1), Z is Z1*X.

```

Debugging

trace - step-by-step view of execution

CALL: is about to invoke a new goal/ match a new clause

RETRY: is doing backtracking(match another clause)

EXIT: a predicate call is successful, and is returning

FAIL: a goal cannot be satisfied, backtracking goes to another part of tree

```
forestDwelling(Animal) :-
    habitat(Animal,_), % add this line
    \+ tropical(Animal). %uses habitat
% Assign Animal to a habitat to work with variables
```

```
is_dog(lab). is_dog(poodle). is_dog(bulldog).
animal(monkey). animal(guinea_pig). animal(lab).
```

```
count_dogs([ListHead|ListTail],Count):-
    is_dog(ListHead), NewCount is Count+1,
    count_dogs(ListTail,NewCount).
count_dogs([ListHead|ListTail],Count):-
    \+ListHead, count_dogs(ListTail,Count).
% Doesn't actually *return* the final count
```

```
non_dog(A):-
    \+is_dog(A),
    animal(A).
% Doesn't get past \+is_dog term
```

```
grandfather(Gramps,Kid):-
    parent(Gramps,parent),
    parent(Parent,Kid).
% Note the 'parent' instead of 'Parent'
```

```
double([], []).
double([N | T], [N2|T2]) :-
    N2 is N*2,
    double(T, Soln2).
% Doubles all elements in list, but Soln2\=T2
```

```
% Move head of one list to head of another
movehead([ListAHead|ListATail],ListB,NewList):-
    append(ListAHead,ListB,NewList).
% ListAHead should be [ListAHead]
```

Stacks & Queues

```
%push(Item,Stack,NewStack)
push(Item,Stack,[Item|Stack]).
```

```
%pop(Item,Stack,NewStack)
pop(Item,[Item|Rest],Rest).
```

```
%top(Top,Stack). Top = top_of_stack
top(Top,[Top|_]).
```

```
is_empty([]).
```

```
%Queue
```

```
shove(Item,Queue,New):-
    append(Queue,[Item],New).
yank(Item,[Item|Stack],Stack).
```

Built-in Functions

```
random(X), write(X), nl,
read(X) ← reads next term from input stream,
get(X) ← read single character,
tab(X) ← print X spaces,
put(X) ← write single character,
see(X) ← opens input stream to come from file X (default
'user'),
seeing(X) ← indicates current input stream,
seen ← closes input stream, resets input as user,
tell(X) ← opens file X as target for output stream,
telling(X) ← where you are writing to,
told ← closes output stream, resets to user,
sort(L,S) ← sorts the list, but not nested lists,
setof(T,G,Set) ← unique elements and sorted
```

Cuts (!)

1. All the clauses before the first clause with a cut are executed with normal backtracking.
2. If the goals before the cut never succeed, the cut does not activate, and the subsequent clause is used, as normal.
3. If the goals before the cut succeed, the cut activates:
 - a) backtracking back to goals before the cut cannot occur
 - b) backtracking to subsequent clauses after the one with the cut cannot occur → that clause with the activated cut is committed
 - c) the goals after the cut are executed with normal backtracking

Testing and Manipulation

```
var(X) ← succeeds if argument X is an uninstantiated variable,
nonvar(X) ← succeeds if argument X is instantiated to something,
atom(X) ← succeeds if X is a non-numeric constant,
integer(X) ← succeeds if X is an integer,
ground(X) ← succeeds if X is instantiated to something that has no uninstantiated variables eg. ground(s(1,X)) fails, but ground(s(1,2)) succeeds,
functor(T, F, N) ← succeeds if term T has functor name F and arity N,
name(A, X) ← converts atom A into list of ascii values (integers); works backwards too
```

Equality

```
= unifies (eg: A = cat → true)
== does not unify (eg: A == cat → false) they have to be identically equal
```

Binary Tree

Binary trees are very useful when they are sorted: keys on the left branch are less than the node, and keys on right are greater

```
% add_bintree(Key, OldTree, NewTree)...
add_bintree(Key, nil, tree(nil, Key, nil)).
add_bintree(Key, tree(L, Key, R), tree(L, Key, R)).
add_bintree(Key, tree(L, K, R), tree(L2, K, R)) :-
    Key < K,
    add_bintree(Key, L, L2).
add_bintree(Key, tree(L, K, R), tree(L, K, R2)) :-
    Key > K,
    add_bintree(Key, R, R2).
```

Definite Clause Grammar (DCG)

```
nonterminal→nonterminals
nonterminal→terminals
nonterminal→nonterminals & terminals
Terminals terminate when called. Eg) letter→[a]
```

Example 1:

```
s-->[a],[b].
s-->[a],s,[b].

?- s(X,[]). X = [a, b] ; X = [a, a, b, b] ; etc

?- s([a],[]). false.

?- s([a,b],[]). true.

?- s([a,b,b],[]). false.

?- s([a,b,b],X). X = [b].
% [a,b] is consumed and [b] is left over

?- s([a,b],X). X = [].

?- s([],X). false. % add s-->[] to be true
```

Example 2:

```
sentence-->noun_phrase(subject), verb_phrase.
verb_phrase-->verb, noun_phrase(object).
noun_phrase(_)-->determiner,noun. % eg: 'the cat'
noun_phrase(X)-->pronoun(X). % eg: 'he', 'him'

determiner-->[a]. determiner-->[the].
noun-->[cat]. noun-->[mouse].
verb-->[scares]. verb-->[hates].
pronoun(subject)-->[he]. pronoun(subject)-->[she].
pronoun(object)-->[him]. pronoun(object)-->[her].
OR
determiner-->[Word],{lex(Word,det)}.
noun-->[Word],{lex(Word,noun)}. % Word = noun
verb-->[Word],{lex(Word,verb)}.
pronoun(subject)-->[Word],{lex(Word,pro_sub)}.
pronoun(object)-->[Word],{lex(Word,pro_ob)}.
lex(a,det). lex(the,det).
lex(cat,noun). lex(mouse,noun).
lex(scars,verb). lex(hates,verb).
lex(he,pro_sub). lex(she,pro_sub).
lex(him,pro_ob). lex(her,pro_ob).
```

Constraint Logic Programming (CLP)

CLP (or DCG) question has something evil in it!

Example 1: List the possible values for X, Y, and Z

Note: \backslash is a unification of domains. Thus, $1.3\backslash 5.7$ means "1 to 3, or 5 to 7".

```
[X,Y,Z] ins 1..3\5\7..9,
Y#=X+1,
Z#=Y+2.
```

```
X: 2
Y: 3
Z: 5
```

```
[X,Y,Z] ins 1..3\5\7..9,
Y#=X+1,
Z#>Y.
```

```
X: 1..2\7
Y: 2..3\8
Z: 3\5\7..9
```

Example 2: Calculate temperature

```
:- use_module(library(clpr)).
temperature(Celsius,Fahrenheit):-
    {Celsius*9/5+32=Fahrenheit}.
```

Assignment 3

Question 1: Tokenizer

```
?- tokenize.
|: hi there
[hi, there]
```

```
tokenize:-
    read_line_to_codes(user_input,String),
    parse(String,[],[]).
```

```
parse([],ListOfTokens,CurrentToken):- % base case
    reverse(CurrentToken,CompleteToken),
    name(Token, CompleteToken), % convert to chars
    Tmp = [Token|ListOfTokens], % add to list
    reverse(Tmp,Output), % reverse complete list
    write(Output), !.
```

```
parse([H|T], ListOfTokens, CurrentToken):-
    H \= 32, % add char to current token
    parse(T, ListOfTokens, [H|CurrentToken]);
    H = 32, % else add current token to list
    reverse(CurrentToken,CompleteToken),
    name(Token, CompleteToken),
    parse(T, [Token|ListOfTokens], []).
    % insert token into list, and reset current token
```

```
reverse(List,Reversed):- rev(List,[],Reversed).
rev([],Rev,Rev).
rev([A|T],L,Rev):- rev(T,[A|L],Rev).
```

```
% This version uses a built-in predicate
tokenizer(String):-
    concat_atom(Output,' ',String), write(Output).
```

Question 2: Postfix calculator

Uses the same tokenizer as Q1, but instead of
 “reverse(Tmp,Output), write(Output), !.”
 It has “reverse(Tmp,List), !, calculate(List, []).” and...

```
push(Item,Stack,[Item|Stack]).
pop(Item,[Item|Rest],Rest).
top(Top,[Top|_]).
```

% Examples:

```
% ?- push(2,[],Stack), push(3,Stack,NewStack).
% Stack = [2],
% NewStack = [3, 2].
```

```
% ?- pop(H, [a,b,c],Stack).
% H = a,
% Stack = [b, c].
```

```
calculate([],Stack):- % base case
    top(Top, Stack), write('Solution: '), write(Top).
calculate([H|T], Stack):-
    % First check for operators
    H = +, % ADD the top 2 values
    pop(Operand2, Stack, NewStack), % pop 2 values
    pop(Operand1, NewStack, NewerStack),
    Result is Operand1 + Operand2, % apply operator
    push(Result, NewerStack, NewestStack), !,
    calculate(T, NewestStack); % continue
    .
    . similar for Subtract, Multiply, Divide
    .
    % Else push the operand onto the stack
    push(H, Stack, NewStack),
    calculate(T, NewStack).
```

Question 3: 0-1 Knapsack

```
knapsack(Capacity):-
    abolish(sack/1),
    asserta(sack([])), % contents of best sack
    abolish(capacity/1),
    asserta(capacity(Capacity)),
    abolish(maxW/1),
    asserta(maxW(0)), % best weight
    abolish(maxV/1),
    asserta(maxV(0)), % best profit
    abolish(item/3), % cleanup any traces
    readFile(input), % Linux
    writeln('Data loaded.'),
    writeln('Finding optimal loot...'),nl,
    bruteForce,
    maxV(V), maxW(W), sack(S),
    writeln(V), writeln(W), writeln(S),!.
```

```
readFile(X):-
    see(X), seeing(InStream),nl,
    repeat, read_line_to_codes(InStream, String),
    tokenize(String),
    String = end_of_file, !, seen.
```

```
tokenize(end_of_file). % terminal condition
```

```
tokenize(String):-
    String \= end_of_file, % don't tokenize eof
    parse(String,[],[]).
```

```
% parse and reverse are the same, except
% addToKB(Output) added before cut, delimiter is 9
```

```
addToKB([Object, Weight, Value |_]):-
    assert(item(Object,Weight,Value)),
    write(Object),
    write(': W='), write(Weight),
    write(', V='), write(Value),nl.
```

```
count([],0).
count([_|T],N) :- count(T,Sum),N is Sum+1.
```

```
perm(List,[H|Perm]):-
    delete(H,List,Rest), perm(Rest,Perm).
perm([],[]).
```

```
delete(X,[X|T],T).
delete(X,[H|T],[H|NT]):- delete(X,T,NT).
```

```
/* Create all variations */
modified_varia(0,_).
modified_varia(N,X):-
    Z is N-1,
    bagof(L, varia(N,X,L), LL),
    bruteForce(LL),
    modified_varia(Z,X).
```

```
varia(0,[],[]).
varia(N,L,[H|Varia]):-
    N > 0, N1 is N-1,
    delete(H,L,Rest), varia(N1,Rest,Varia).
```

```
bruteForce:-
    findall([X,Y,Z],item(X,Y,Z),Items),
    count(Items, N), modified_varia(N, Items).
```

```
/* Evaluate every variation to find the optimal loot */
bruteForce([]).
```

```
bruteForce([H|T]):-
    computeLoot(H,0,0,H),
    bruteForce(T).
```

```
/* Compute the current loot variation */
computeLoot([], TotalW, TotalV, Copy):-
    maxV(MaxV),
```


palindrome:-	69 E
read_line_to_codes(user,String),	70 F
palindrome(String,[]).	71 G
	72 H
% Check for invalid characters	73 I
palindrome-->[32],palindrome.	74 J
palindrome-->[Char],palindrome,[Char],[32].	75 K
palindrome-->[_],[32].	76 L
	77 M
palindrome-->[39],palindrome.	78 N
palindrome-->[Char],palindrome,[Char],[39].	79 O
palindrome-->[_],[39].	80 P
	81 Q
% There are 3 cases of what a palindrome is	82 R
palindrome-->[Char],palindrome,[Char].	83 S
palindrome-->[_]. % single letter	84 T
palindrome-->[]. % an empty string	85 U
	86 V
<u>ASCII Table</u>	87 W
	88 X
32 (white space)	89 Y
33 !	90 Z
34 "	91 [
35 #	92 \
36 \$	93]
37 %	94 ^
38 &	95 _
39 ' ,	96 `
40 (97 a
41)	98 b
42 *	99 c
43 +	100 d
44 ,	101 e
45 -	102 f
46 .	103 g
47 /	104 h
48 0	105 i
49 1	106 j
50 2	107 k
51 3	108 l
52 4	109 m
53 5	110 n
54 6	111 o
55 7	112 p
56 8	113 q
57 9	114 r
58 :	115 s
59 ;	116 t
60 <	117 u
61 =	118 v
62 >	119 w
63 ?	120 x
64 @	121 y
65 A	122 z
66 B	123 {
67 C	124
68 D	

125 }

126 ~

Arbitrary

- Procedural programming is "how to" programming:
Understand program in terms of goals to solve and in which order to solve them.

- Declarative programming is "what to" programming:
Understand each program predicate at a high level.

- Backtracking causes execution to revert back to last place a clause was successful, and move on to the next one.

- Prolog is made up of (1) Facts, (2) Queries, (3) Rules

- Green cuts: prune tree w/o affecting final answer

- Red cuts: prune tree and affect final answer

- if-then-else* can be emulated via a cut:

```
if T is true, do Q, else R
```

```
P :- (T -> Q ; R).
```

```
P :- T, !, Q.
```

```
P :- R.
```

- Write all matches without hitting next (;)

```
likes(boy,girl). likes(programmer,code).
```

```
write_likes:- likes(X,Y), write(X),
```

```
    write(' likes '), writeln(Y), fail.
```

```
write_likes. % finish as true
```